

# STABILITY OF COMMUNICATIONS NETWORKS IN THE PRESENCE OF DELAYS

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Abstract: In this paper, we present a study of the destabilizing effects of delay in communications networks. We apply results from a companion paper to find the range of delays for which a Smith predictor controller guarantees the stability of an ATM/ABR network.

Keywords: Time Delay, Stability, Smith Predictor

## 1. THE ATM/ABR EXAMPLE

Congestion control in the Available Bit Rate (ABR) class of Asynchronous Transfer Mode (ATM) networks poses interesting challenges due to the presence of multiple-delays, magnitude and rate constraints on the inputs, amplitude limitation on the state, and additive disturbances. The ABR class is designed as a best-effort class for applications such as file transfer and email. Thus, no service guarantees are required, but the source of data packets controls its data rate, using a feedback signal provided by switches downstream that measure the congestion of the network. Due to the presence of this feedback, many classical and advanced control theory concepts have been suggested to deal with the congestion control problem in the ATM/ABR case.

We will focus our modeling on one particular queue in the network which is associated with a link shared by many virtual circuits D. Cavendish and Gerla (1996). A virtual circuit is established for any two stations wishing to communicate, by informing all switches between them of their

requirements. Assuming that the flow of packets is conserved, the queue level model for each buffer in the ATM network (as proposed in D. Cavendish and Gerla (1996)) is

$$\dot{x}(t) = \sum_{i=1}^n u_i(t - T_i) - d(t) \quad (1)$$

where

- $x(t)$  is the queue level associated with the considered link;
- $n$  is the number of *virtual circuits* sharing the queue associated with the considered link;
- $u_i(t)$  is the inflow cell rate caused by the  $i$ th *virtual circuit*;
- $T_i$  is the propagation delay from the  $i$ th source to the queue, and is usually uncertain;
- $d(t)$  is the rate of packets leaving the queue.

More details on this model and its physical interpretation is available from Mascolo (2000) (see also Nicolescu (2001) p. 60-63). Equation (1) appears to be linear since the saturation effect due to the limited buffer capacity has been neglected. Our control design however, must ensure that this condition is actually satisfied, so that the controller is not saturated. The author in Mascolo (2000) proposes the Smith predictor control

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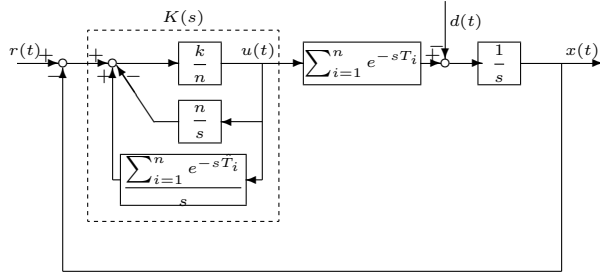


Fig. 1. The controlled system.

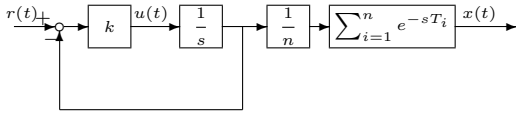


Fig. 2. The desired closed-loop system.

scheme of Figure 1. The objectives of the control law proposed in Mascolo (2000) are to guarantee

- *Stability*<sup>2</sup>:

$$x(t) \leq r^0 \quad t \geq 0$$

where  $r^0$  is the queue capacity. This condition guarantees no cell loss and is not the usual stability requirement;

- *Full Link Utilization*:

$$x(t) > 0 \quad t > T_{tr} \quad (2)$$

The time  $T_{tr}$  in (2) mainly accounts for the transient time of the dynamics. Let  $\hat{T}_i$  denote the estimated values of the delays  $T_i$ . Assuming that  $\hat{T}_i = T_i$  and disregarding for the moment the presence of disturbance  $d(t)$ , the scheme in Figure 1 can be shown to be equivalent to the one in Figure 2 Mascolo (2000). In this case, the designer knows exactly the value of all delays  $T_i; i = 1, \dots, n$ , and the closed-loop system exhibits the following nice properties.

- With  $r(t) = r^0 \cdot 1(t)$  ( $1(t)$  is the unit step) denoting the desired reference level of the queue, and  $d(t) = a \cdot 1(t)$ , with  $0 \leq a \leq 1$  the *stability property* (boundedness) is satisfied.
- Further, with  $d(t) = a \cdot 1(t)$ , with  $0 \leq a \leq 1$  the *full utilization property* is satisfied, provided that Mascolo (2000)

$$r^0 > a \left( \frac{1}{k} + \frac{1}{n} \sum_{i=1}^n T_i \right).$$

The technique which is used in order to obtain the desired closed-loop system is the well-known Smith's principle Smith (1957) Power and Simpson (1978).

The main drawback of the approach of Mascolo (2000) is that it assumes that the propagation

delays  $T_i$  are *exactly* known. When this is not the case, even stability, *in the sense defined above*<sup>3</sup>, can be lost. In order to illustrate this limitation and the potential of losing stability (boundedness), a simulation was performed. Assuming that  $n = 4$ ,  $T_1 = 10$  sec,  $T_2 = 30$  sec,  $T_3 = 60$  sec,  $T_4 = 120$  sec,  $k = 0.1$ , and  $r^0 = 40$ , the stability condition is satisfied (see Figure 3) if  $\hat{T}_i = T_i$   $i = 1, \dots, 4$ .

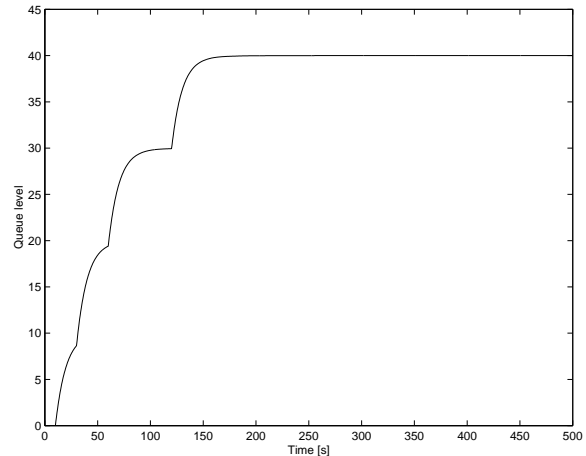


Fig. 3. Queue level when the propagation delays are assumed to be known.

Next, the actual delays  $T_i; i = 1, \dots, n$  are perturbed by 5%, without modifying their assumed values  $\hat{T}_i; i = 1, \dots, n$ . As shown in Figure 4, stability is lost since  $x(t)$  is greater than  $r^0$  during some time intervals. Stability is regained however if the controller gain  $k$  is changed to  $k = 0.01$  (see Figure 5) at the expense of a slower response. When the  $T_i$ s however are not known, the designer cannot apriori design a Smith-predictor-type controller in order to guarantee the desired degree of stability and performance. Moreover, even BIBO stability may be lost if one does not take the effects of the uncertainty in the delays.

## 2. DELAYS IN THE ATM MODEL

The transfer function of the ATM system is

$$X(s) = \frac{\frac{k}{n} \sum_{i=1}^n e^{-sT_i}}{s + k - \frac{k}{n} \left( \sum_{i=1}^n e^{-sT_i} - e^{-s\hat{T}_i} \right)} R(s) - \frac{s + \frac{k}{n} (n - \sum_{i=1}^n e^{-s\hat{T}_i})}{s \left( s + k - \frac{k}{n} \left( \sum_{i=1}^n e^{-sT_i} - e^{-s\hat{T}_i} \right) \right)} d$$

First of note that the transfer function

$$\frac{s + \frac{k}{n} (n - \sum_{i=1}^n e^{-s\hat{T}_i})}{s}$$

<sup>2</sup> The network community refers to this as a “stability” condition. Stability in the dynamic systems sense will guarantee BIBO stability and then care must be taken to ensure the bound  $r^0$  is not violated.

<sup>3</sup> That is,  $x(t) \leq r^0$  is violated.

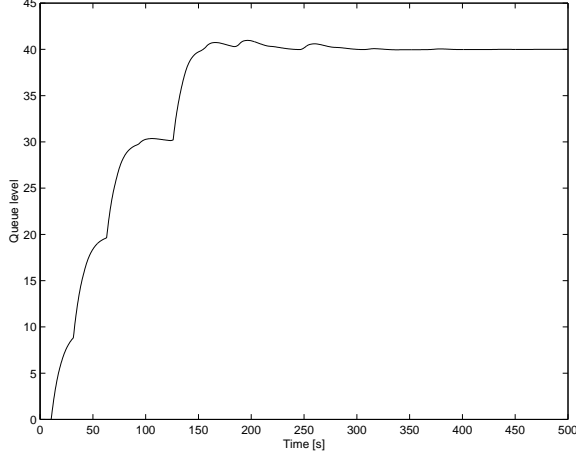


Fig. 4. Queue level when the propagation delays are perturbed by 5%.

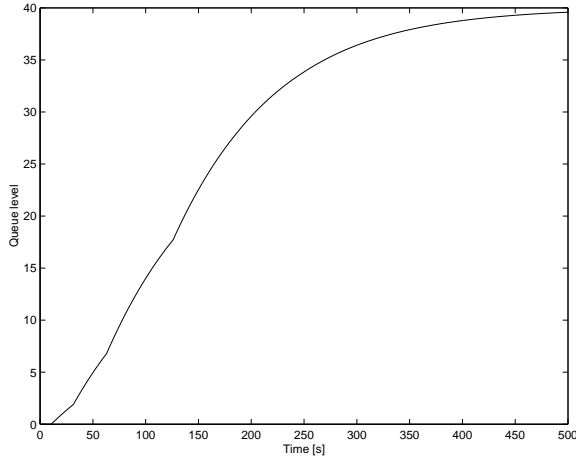


Fig. 5. Queue level when the propagation delays are perturbed by 5% and  $k = 0.01$ .

is stable since the only possible pole location is  $s = 0$  and there it follows that

$$\begin{aligned} & \left. \frac{s + \frac{k}{n}(n - \sum_{i=1}^n e^{-s\hat{T}_i})}{s} \right|_{s=0} \\ &= \lim_{s \rightarrow 0} \frac{s + \frac{k}{n}(n - \sum_{i=1}^n e^{-s\hat{T}_i})}{s} \\ &= \lim_{s \rightarrow 0} \frac{s + \frac{k}{n}(n \sum_{i=1}^n \hat{T}_i)s}{s} \\ &= 1 + k \sum_{i=1}^n \hat{T}_i \neq \infty. \end{aligned}$$

Consequently, the stability of the ATM system is reduced to studying the stability of  $s + k - \frac{k}{n}(\sum_{i=1}^n e^{-sT_i} - e^{-s\hat{T}_i})$ . In this section, being stable means

$$s + k - \frac{k}{n}(\sum_{i=1}^n e^{-sT_i} - e^{-s\hat{T}_i}) \neq 0 \text{ for } \text{Re}(s) \geq 0 \quad (3)$$

In order to analyze the stability of this system, some further refinements in the model are now made. Let  $T_i = m_i h_0$ ,  $\hat{T}_i = m_i h$  for

$i = 1, \dots, n$  where the  $m_i$  are positive integers<sup>4</sup>. Define  $c(z) \triangleq \sum_{i=1}^n z^{m_i}$  so that  $c(e^{-sh_0}) = \sum_{i=1}^n e^{-sT_i}$ ,  $c(e^{-sh}) = \sum_{i=1}^n e^{-s\hat{T}_i}$ . The characteristic equation of the ATM system (3) may be rewritten as

$$a(s, e^{-h_0s}, e^{-hs}) \triangleq a_0(s) + a_1(s) [c(e^{-sh_0}) - c(e^{-sh})] \quad (4)$$

where  $a_0(s) = s + k$ ,  $a_1(s) = -k/n$ . Here the delays  $h_0, h$  are noncommensurate in general. By design, the compensator  $G_c(s) = k$  is chosen so that (4) is stable for  $h = h_0$ , that is,  $a_0(s)$  is stable. The interest here is to determine the range of values of  $h$  about  $h_0$  for which (4) remains stable. As  $h$  is decreased (or increased) in value from  $h_0$ , there is a first value  $h^*$  for which the roots of (4) are on the  $j\omega$  axis, that is,

$$\begin{aligned} & a(s, e^{-h_0s}, e^{-h^*s}) \\ & \triangleq a_0(s) + a_1(s) (c(e^{-sh_0}) - c(e^{-sh^*})) \\ & = 0 \end{aligned} \quad (5)$$

for some  $s^* = j\omega^*$ .

To exploit this idea ( $z_0 = e^{-sh_0}$ ,  $z = e^{-sh}$ ), define ( $m = \max_i \{m_i\}$ )

$$\begin{aligned} & a(s, z_0, z) \\ &= a_0(s) + a_1(s) (c(z_0) - c(z)), \quad c(z) = \sum_{i=1}^n z^{m_i} \\ & \tilde{a}(s, z_0, z) \\ & \triangleq z_0^m z^m a(-s, 1/z_0, 1/z) \\ &= z_0^m z^m [a_0(-s) + a_1(-s) (c(1/z_0) - c(1/z))] \\ &= z_0^m z^m a_0(-s) + a_1(-s) (z^m c(z_0) - z_0^m c(z)) \end{aligned}$$

Solving  $a(s, z_0, z) = 0$  to get  $c(z) = (a_0(s) + a_1(s)c(z_0))/a_1(s)$  and substituting this into  $\tilde{a}(s, z_0, z) = 0$  gives

$$\begin{aligned} 0 &= z_0^m z^m a_0(-s) + \\ & a_1(-s) \left( z^m c(z_0) - z_0^m \frac{a_0(s) + a_1(s)c(z_0)}{a_1(s)} \right) \\ z^m &= \frac{z_0^m (a_1(-s)a_0(s) + a_1(s)a_1(-s)c(z_0))}{(z_0^m a_1(s)a_0(-s) + a_1(-s)a_1(s)c(z_0))} \\ &= \frac{a_1(-s)a_0(s) + a_1(s)a_1(-s)c(z_0)}{a_1(s)a_0(-s) + a_1(-s)a_1(s)c(1/z_0)} \end{aligned}$$

where  $z_0^{-m}c(z_0) \equiv c(1/z_0)$ . With  $s = j\omega$ ,  $z = e^{-jh\omega}$ ,  $z_0 = e^{-jh_0\omega}$  this becomes

<sup>4</sup> For all practical purposes, there will always be  $h_0, h$  such that the delays can be written as integer multiples of these values.

$$z^m = e^{-jm h \omega}$$

$$\begin{aligned} &= \frac{a_1(-j\omega)a_0(j\omega) + a_1(j\omega)a_1(-j\omega)c(e^{-jh_0\omega})}{a_1(j\omega)a_0(-j\omega) + a_1(-j\omega)a_1(j\omega)c(e^{jh_0\omega})} \\ &= \left| \frac{a_1(-j\omega)a_0(j\omega) + a_1(j\omega)a_1(-j\omega)c(e^{-jh_0\omega})}{a_1(j\omega)a_0(-j\omega) + a_1(-j\omega)a_1(j\omega)c(1/e^{-jh_0\omega})} \right| \\ &\times e^{j\angle \frac{a_1(-j\omega)a_0(j\omega) + a_1(j\omega)a_1(-j\omega)c(e^{-jh_0\omega})}{a_1(j\omega)a_0(-j\omega) + a_1(-j\omega)a_1(j\omega)c(e^{jh_0\omega})}} \\ &= e^{2j\angle(a_1(-j\omega)a_0(j\omega) + a_1(j\omega)a_1(-j\omega)c(e^{-jh_0\omega}))} \end{aligned}$$

and solving for  $z$

$$\begin{aligned} z &= e^{-jh\omega} \\ &= e^{j \frac{2\angle(a_1(-j\omega)a_0(j\omega) + a_1(j\omega)a_1(-j\omega)c(e^{-jh_0\omega})) + 2\pi k}{m}} \quad (6) \end{aligned}$$

for  $k = 0, 1, \dots, m-1$ . By definition, for  $h = h^*$  corresponds to  $a(s, e^{-h_0j\omega}, e^{-h^*j\omega}) = 0$  for some  $s^* = j\omega^*$  where  $e^{-h^*j\omega}$  is given by the right-hand side of (6) for  $k = 0, 1, \dots, m-1$ . Thus, the right-hand side of (6) is substituted into (5) to get

$$0 = a_0(j\omega) + a_1(j\omega) \left[ c(e^{-jh_0\omega}) - c(e^{-jh^*\omega}) \right] \Big|_{e^{-jh^*\omega} = \text{rhs of (6)}}$$

For each  $k = 0, 1, \dots, m-1$  this equation is solved for the corresponding  $\omega_{ki}$ ,  $i = 1, 2, \dots$  which are then substituted back into (6) so for  $k = 0, 1, \dots, m-1$  it follows that

$$\begin{aligned} e^{-jh^*\omega_{ki}} &= \\ e^{j(2\angle(a_1(-j\omega_{ki})a_0(j\omega_{ki}) + a_1(j\omega_{ki})a_1(-j\omega_{ki})c(e^{-jh_0\omega_{ki}})) + 2\pi k)/n} \end{aligned}$$

Finally, for each  $\omega_{ki}$  this is solved  $h_{ki}^*$ . Then

$$\begin{aligned} h_1^* &= \max_{k,i} \{h_{ki}^* : h_{ki}^* < h_0\} \\ h_2^* &= \min_{k,i} \{h_{ki}^* : h_{ki}^* > h_0\}. \end{aligned}$$

Finally, for  $h_1^* < h < h_2^*$

$$\begin{aligned} a(s, e^{-h_0s}, e^{-hs}) &\triangleq \\ a_0(s) + a_1(s) [c(e^{-sh_0}) - c(e^{-sh})] &\neq 0 \text{ for } \text{Re}\{s\} \geq 0 \end{aligned}$$

and

$$\begin{aligned} a(s, e^{-h_0s}, e^{-h_1^*s}) &= 0 \text{ for some } s = j\omega \\ a(s, e^{-h_0s}, e^{-h_2^*s}) &= 0 \text{ for some } s = j\omega \end{aligned}$$

### 3. CONCLUSIONS

In this paper, we have applied results from our companion paper Chiasson and Abdallah (2001) in order to exactly determine the range of delays for which Smith predictor controllers remain stabilizing. The results are specialized to ATM/ABR networks, but can also be applied in many areas of queuing theory and communications.

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